

Triangular Decomposition of Tensor Product of Simple Graphs

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Abstract: Let $G = (V, E)$ be a simple connected graph of order p and size q . If $\{G_1, G_2, \dots, G_n\}$ are edge disjoint subgraphs of G such that $E(G) = E(G_1) \cup E(G_2) \cup \dots \cup E(G_n)$ then $\{G_1, G_2, \dots, G_n\}$ is said to be a Decomposition of a graph G . A graph of size $q = \binom{n+2}{3}$ is said to have a Triangular decomposition (TD) if G can be decomposed into n - subgraphs $\{G_1, G_2, \dots, G_n\}$ such that each subgraphs G_i is connected and $|E(G_i)| = \binom{i+1}{2}$ for $1 \leq i \leq n$. In this paper we investigate Triangular decomposition of Tensor product of simple graphs.

Keywords: Triangular decomposition, Tensor product, Wheel Graph.

1. Introduction

A graph G , referred to here is an undirected connected graph without loops or multiple edges. Let $G_1 = (V_1, E_1)$ and $G_2 = (V_2, E_2)$ be two graphs. The Tensor product $G = G_1 \wedge G_2$ is defined as a graph with vertex set $V_1 \times V_2$. Edge set is defined as follows: If $w_1 = (u_1, v_1)$ and $w_2 = (u_2, v_2)$ are two vertices of G with $u_i \in V_1$ and $v_i \in V_2$ ($i=1,2$) then $w_1 w_2 \in E(G)$ if and only if $u_1 u_2 \in E_1$ and $v_1 v_2 \in E_2$. [4] The concept of Continuous Monotonic Decomposition of Graph was introduced by N.Gnana Dhas and J.Paulraj Joseph. [6] S.Asha, and R.kala discussed on Continuous Monotonic Decomposition of some special class of Graphs. Terms not defined here are used in the sense of Harary [1]. In this paper we proved some results on Triangular decomposition of Tensor product of simple graphs.

Definition 1.1 A decomposition of a graph G is a collection of edge disjoint subgraphs $\{G_1, G_2, G_3, \dots, G_n\}$ of G such that every edge of G belongs to exactly one of the subgraph G_i .

Definition 1.2 A graph G of size $q = \binom{n+2}{3}$ is said to have a Triangular decomposition (TD) if G can be decomposed into n - subgraphs $\{G_1, G_2, G_3, \dots, G_n\}$ such that each G_i is connected and $|E(G_i)| = \binom{i+1}{2}$ for $1 \leq i \leq n$.

Definition 1.3. The tensor product $G \times H$ of graphs G and H is a graph such that the vertex set of $G \times H$ is the Cartesian product $V(G) \times V(H)$; and vertices (g, h) and (g', h') are adjacent in

$G \times H$ if and only if g is adjacent to g' in G and h is adjacent to h' in H .

2. Triangular Decomposition of Tensor Product of Simple Graphs

Lemma 2.1. If $k = 4s$, then every graph of size $q = \frac{k(k+1)(k+2)}{6}$ can be decomposed into $\{G_1, G_2, G_3, \dots, G_k\}$.

Proof. We have $k = 4s$, $s \in \mathbb{N}$. We prove this theorem by using induction method. When $s = 1$, $k = 4$. Then $q = \frac{k(k+1)(k+2)}{6} = \frac{4 \times 5 \times 6}{6} = 20$ can be decomposed into $\{G_1, G_2, G_3, G_4\}$. Hence the result is true for $s = 1$. Assume that the result is true for $s-1$. Then $k = 4(s-1) = 4s-4$ and $q = \frac{(4s-4)(4s-3)(4s-2)}{6}$ can be decomposed into $\{G_1, G_2, G_3, \dots, G_{4s-4}\}$. Now to prove the result is true for s . Then $k = 4s$ and $q = \frac{4s(4s+1)(4s+2)}{6}$. We have to prove that $\frac{4s(4s+1)(4s+2)}{6}$ can be decomposed into $\{G_1, G_2, G_3, \dots, G_{4s}\}$.

$$\begin{aligned} \text{Now } q &= \frac{4s(4s+1)(4s+2)}{6} \\ &= \frac{64s^3 + 48s^2 + 8s}{6} \\ &= \frac{(4s-4)(16s^2 - 20s + 6)}{6} + \left\{ \frac{(4s-2)(8s-4)}{2} + \frac{(4s)(8s)}{2} \right\} \\ &= \frac{(4s-4)(16s^2 - 8s - 12s + 6)}{6} + \left\{ \frac{(4s-2)(4s-3+4s-1)}{2} + \frac{(4s)(4s-1+4s+1)}{2} \right\} \\ &= \frac{(4s-4)(4s-3)(4s-2)}{6} + \left\{ \frac{(4s-3)(4s-2)}{2} + \frac{(4s-2)(4s-1)}{2} + \frac{(4s-1)(4s)}{2} + \frac{4s(4s+1)}{2} \right\}. \end{aligned}$$

Therefore $q = \frac{4s(4s+1)(4s+2)}{6}$ can be decomposed into $\{G_1, G_2, G_3, \dots, G_{4s}\}$. Hence by induction hypothesis if $k = 4s$, then every graph of size $q = \frac{k(k+1)(k+2)}{6}$ can be decomposed into $\{G_1, G_2, G_3, \dots, G_k\}$. This completes the proof.

Lemma 2.2. If $k+1=4s$, then every graph of size $q = \frac{k(k+1)(k+2)}{6}$ can be decomposed into $\{G_1, G_2, G_3, \dots, G_k\}$.

Proof. We have $k = 4s - 1$, $s \in \mathbb{N}$. We prove this theorem by using induction method. When $s = 1$, $k = 3$. Then $q = \frac{k(k+1)(k+2)}{6} = \frac{3 \times 4 \times 5}{6} = 10$ can be decomposed into $\{G_1, G_2, G_3\}$. Hence the result is true for $s = 1$. Assume that the result is true for $s-1$. Then $k = 4(s-1)-1 = 4s-5$ and $q = \frac{(4s-5)(4s-4)(4s-3)}{6}$ can be decomposed into $\{G_1, G_2, G_3, \dots, G_{4s-5}\}$. Now to prove the result is true for s . Then $k = 4s-1$ and $q = \frac{(4s-1)(4s)(4s+1)}{6}$. We have to prove that $q = \frac{(4s-1)(4s)(4s+1)}{6}$ can be decomposed into $\{G_1, G_2, G_3, \dots, G_{4s-1}\}$.

$$\begin{aligned} \text{Now } q &= \frac{(4s-1)(4s)(4s+1)}{6} \\ &= \frac{64s^3 - 4s}{6} \\ &= \frac{(4s-5)(16s^2 - 28s + 12)}{6} + \left\{ \frac{(4s-3)(8s-6)}{2} + \frac{(4s-1)(8s-2)}{2} \right\} \\ &= \frac{(4s-5)(16s^2 - 12s - 16s + 12)}{6} + \left\{ \frac{(4s-3)(4s-4+4s-2)}{2} + \frac{(4s-1)(4s-2+4s)}{2} \right\} \\ &= \frac{(4s-5)(4s-4)(4s-3)}{6} + \left\{ \frac{(4s-4)(4s-3)}{2} + \frac{(4s-3)(4s-2)}{2} + \frac{(4s-2)(4s-1)}{2} + \frac{(4s-1)4s}{2} \right\}. \end{aligned}$$

Therefore $q = \frac{(4s-1)(4s)(4s+1)}{6}$ can be decomposed into $\{G_1, G_2, G_3, \dots, G_{4s-1}\}$. Hence by induction hypothesis if $k+1 = 4s$, then every graph of size $q = \frac{k(k+1)(k+2)}{6}$ can be decomposed into $\{G_1, G_2, G_3, \dots, G_k\}$. This completes the proof.

Lemma 2.3. If $k+2 = 4s$, then every graph of size $q = \frac{k(k+1)(k+2)}{6}$ can be decomposed into $\{G_1, G_2, G_3, \dots, G_k\}$

Proof. We have $k = 4s-2$, $s \in \mathbb{N}$. We prove this theorem by using induction method. When $s=1$, $k = 2$. Then $q = \frac{k(k+1)(k+2)}{6} = \frac{2 \times 3 \times 4}{6} = 4$ can be decomposed into $\{G_1, G_2\}$. Hence the result is true for $s = 1$. Assume that the result is true for $s-1$. Then $k = 4(s-1)-2 = 4s-6$ and $q = \frac{(4s-6)(4s-5)(4s-4)}{6}$ can be decomposed into $\{G_1, G_2, G_3, \dots, G_{4s-6}\}$. Now to prove the result is true for s . Then $k = 4s-2$ and $q = \frac{(4s-2)(4s-1)(4s)}{6}$. We have to prove that $q = \frac{(4s-2)(4s-1)(4s)}{6}$ can be decomposed into $\{G_1, G_2, G_3, \dots, G_{4s-2}\}$.

$$\begin{aligned} \text{Now } q &= \frac{(4s-2)(4s-1)(4s)}{6} = \frac{64s^3 - 48s^2 + 8s}{6} \\ &= \frac{(4s-6)(16s^2 - 16s - 20s + 20)}{6} + \left\{ \frac{(4s-4)(8s-8)}{2} + \frac{(4s-2)(8s-4)}{2} \right\} \\ &= \frac{(4s-6)(4s-5)(4s-4)}{6} + \left\{ \frac{(4s-5)(4s-4)}{2} + \frac{(4s-4)(4s-3)}{2} + \frac{(4s-3)(4s-2)}{2} + \frac{(4s-2)(4s-1)}{2} \right\} \end{aligned}$$

Therefore $q = \frac{(4s-2)(4s-1)(4s)}{6}$ can be decomposed into $\{G_1, G_2, G_3, \dots, G_{4s-2}\}$. Hence by induction hypothesis if $k+2 = 4s$, then every graph of size $q = \frac{k(k+1)(k+2)}{6}$ can be decomposed into $\{G_1, G_2, G_3, \dots, G_k\}$. This completes the proof.

Theorem 2.4. For any integer m , the Path graph $P_m \wedge K_2$ admits a Triangular Decomposition $\{G_1, G_2, \dots, G_k\}$ if and only if there exists an integer k satisfying the following properties:

(i) $k = 4r$ or $k = 4r-1$ or $k = 4r-2$, $r \in \mathbb{N}$.

(ii) $2m-2 = \frac{k(k+1)(k+2)}{6}$

Proof. Let $G = P_m \wedge K_2$. Then $q(G) = 2m - 2$. Assume G has a Triangular Decomposition. By the definition of Triangular Decomposition, $q(G) = \frac{k(k+1)(k+2)}{6}$.

Hence $2m-2 = \frac{k(k+1)(k+2)}{6}$.

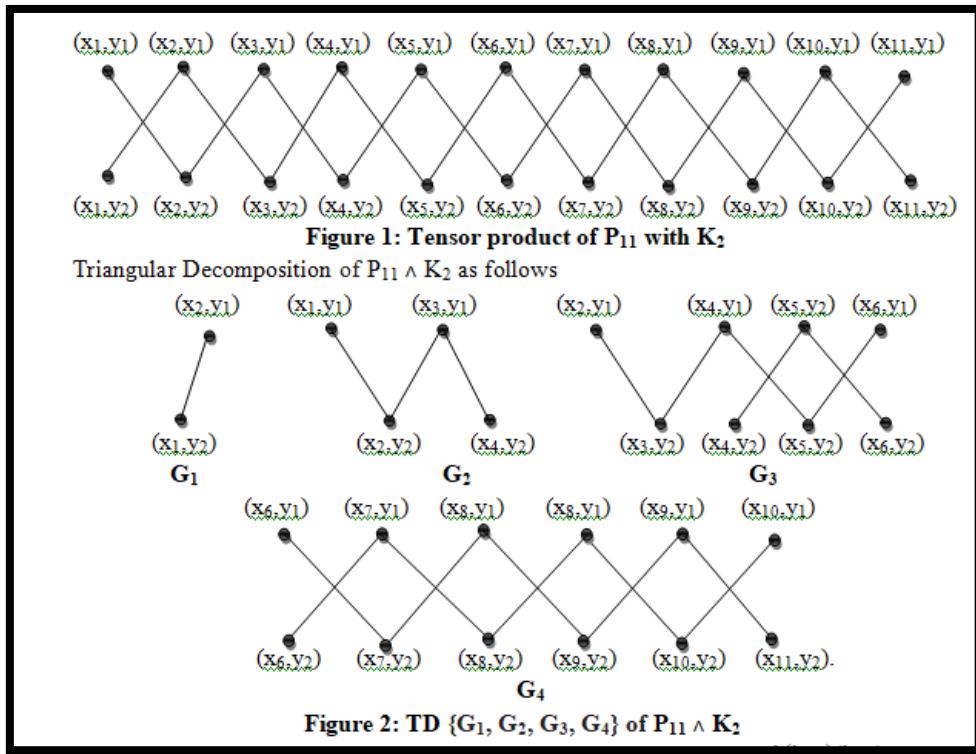
$$\Rightarrow m = \frac{k^3 + 3k^2 + 2k + 12}{12}.$$

Since m is an integer, $k = 4r$ or $k = 4r-1$ or $k = 4r-2$, $r \in \mathbb{N}$.

Conversely assume (i) $k = 4r$ or $k = 4r-1$ or $k = 4r-2$, $r \in \mathbb{N}$. (ii) $2m-2 = \frac{k(k+1)(k+2)}{6}$. Let $G = P_m \wedge K_2$. Then $q(G) = 2m-2$. By lemma 2.1, 2.2 and 2.3, $\frac{k(k+1)(k+2)}{6}$ can be decomposed into $\{G_1, G_2, G_3, \dots, G_k\}$. Thus G admits Triangular Decomposition.

Table 2.5: List of first 10 TD of $P_m \wedge K_2$

m	q(G)	Triangular Decomposition
3	4	G_1, G_2 .
6	10	G_1, G_2, G_3 .
11	20	G_1, G_2, G_3, G_4
29	56	$G_1, G_2, G_3, \dots, G_6$
43	84	$G_1, G_2, G_3, \dots, G_7$
61	120	$G_1, G_2, G_3, \dots, G_8$
111	220	$G_1, G_2, G_3, \dots, G_{10}$
144	286	$G_1, G_2, G_3, \dots, G_{11}$
183	364	$G_1, G_2, G_3, \dots, G_{12}$
281	560	$G_1, G_2, G_3, \dots, G_{14}$

Illustration 2.6

Lemma 2.7 If $k+1=4s$ and $s \equiv 0 \pmod{2}$ then every graph of size $q = \frac{k(k+1)(k+2)}{6}$ can be decomposed into $\{G_1, G_2, G_3, \dots, G_k\}$

Proof. We have $k+1 = 4s$ and $s \equiv 0 \pmod{2}$, $s \in \mathbb{N}$. We prove this theorem by using induction method. When $s = 2$, $k = 7$. Then $q = \frac{k(k+1)(k+2)}{6} = \frac{7 \times 8 \times 9}{6} = 84$ can be decomposed into $\{G_1, G_2, G_3, \dots, G_7\}$. Hence the result is true for $s = 2$. Assume that the result is true for $2s-2$. Then $k = 4(2s-2)-1 = 8s-9$ and $q = \frac{(8s-9)(8s-8)(8s-7)}{6}$ can be decomposed into $\{G_1, G_2, G_3, \dots, G_{8s-9}\}$. Now to prove the result is true for $2s$. Then $k = 8s-1$ and $q = \frac{(8s-1)(8s)(8s+1)}{6}$. We have to prove that $\frac{(8s-1)(8s)(8s+1)}{6}$ can be decomposed into $\{G_1, G_2, G_3, \dots, G_{8s-1}\}$.

$$\begin{aligned}
 \text{Now } q &= \frac{(8s-1)(8s)(8s+1)}{6} \\
 &= \frac{512s^3 - 8s}{6} \\
 &= \frac{(8s-9)(64s^2 - 120s + 56)}{6} + \left\{ \frac{(8s-7)(16s-14)}{2} \right\} + \left\{ \frac{(8s-5)(16s-10)}{2} \right\} + \left\{ \frac{(8s-3)(16s-6)}{2} \right\} + \\
 &\quad \left\{ \frac{(8s-1)(16s-2)}{2} \right\} \\
 &= \frac{(8s-9)(8s-8)(8s-7)}{6} + \left\{ \frac{(8s-8)(8s-7)}{2} + \frac{(8s-7)(8s-6)}{2} + \frac{(8s-6)(8s-5)}{2} + \frac{(8s-5)(8s-4)}{2} \right\} + \\
 &\quad \left\{ \frac{(8s-4)(8s-3)}{2} + \frac{(8s-3)(8s-2)}{2} + \frac{(8s-2)(8s-1)}{2} + \frac{(8s-1)(8s)}{2} \right\}.
 \end{aligned}$$

Therefore $q = \frac{(8s-1)(8s)(8s+1)}{6}$ can be decomposed into $\{G_1, G_2, G_3, \dots, G_{8s-1}\}$. Hence by induction hypothesis if $k+1 = 4s$ and $s \equiv 0 \pmod{2}$ then every graph of size $q = \frac{k(k+1)(k+2)}{6}$ can be decomposed into $\{G_1, G_2, G_3, \dots, G_k\}$. This completes the proof.

Theorem 2.8. For any integer m , the wheel graph $W_{m+1} \wedge K_2$ has a Triangular Decomposition $\{G_1, G_2, G_3, \dots, G_k\}$ iff there exists an integer k satisfying the following properties.

i) $k = 4r$ or $k = 4r-1$, $r \equiv 0 \pmod{2}$ or $k = 4r-2$, $r \in \mathbb{N}$.

ii) $\frac{k(k+1)(k+2)}{6} = 4m$.

Proof. Let $G = W_{m+1} \wedge K_2$. Then $q(G) = 4m$. Assume $W_{m+1} \wedge K_2$ has a Triangular Decomposition. By the definition of Triangular decomposition, $q(G) = \binom{k+2}{3}$.

$$\text{Hence } 4m = \binom{k+2}{3} = \frac{k(k+1)(k+2)}{6}.$$

$$\Rightarrow m = \frac{k(k+1)(k+2)}{24}.$$

Since m is an integer, $k = 4r$ or $k = 4r-1$, $r \equiv 0 \pmod{2}$ or $k = 4r-2$, $r \in \mathbb{N}$.

Conversely assume (i) $k = 4r$ or $k = 4r-1$, $r \equiv 0 \pmod{2}$ or $k = 4r-2$, $r \in \mathbb{N}$

(ii) $\frac{k(k+1)(k+2)}{6} = 4m$. Let $G = W_{m+1} \wedge K_2$. Then $q(G) = 4m$. By lemma 2.1, 2.7 and 2.3, $\frac{k(k+1)(k+2)}{6}$ can be decomposed into $\{G_1, G_2, G_3, \dots, G_k\}$. Thus G admits Triangular Decomposition.

Table 2.9: List of first 10, TD of $W_{m+1} \wedge K_2$.

m	q(G)	Triangular Decomposition
5	20	G_1, G_2, G_3, G_4
14	56	$G_1, G_2, G_3, \dots, G_6$
21	84	$G_1, G_2, G_3, \dots, G_7$
30	120	$G_1, G_2, G_3, \dots, G_8$
55	220	$G_1, G_2, G_3, \dots, G_{10}$
91	364	$G_1, G_2, G_3, \dots, G_{12}$
140	560	$G_1, G_2, G_3, \dots, G_{14}$
170	680	$G_1, G_2, G_3, \dots, G_{15}$
204	816	$G_1, G_2, G_3, \dots, G_{16}$
285	1140	$G_1, G_2, G_3, \dots, G_{18}$

Illustration 2.10

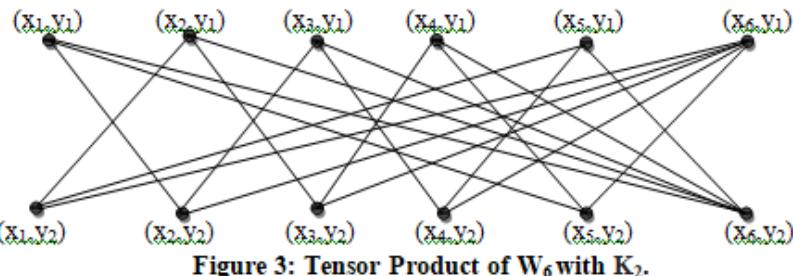


Figure 3: Tensor Product of $W_6 \wedge K_2$.

Triangular Decomposition of $W_6 \wedge K_2$ as follows.

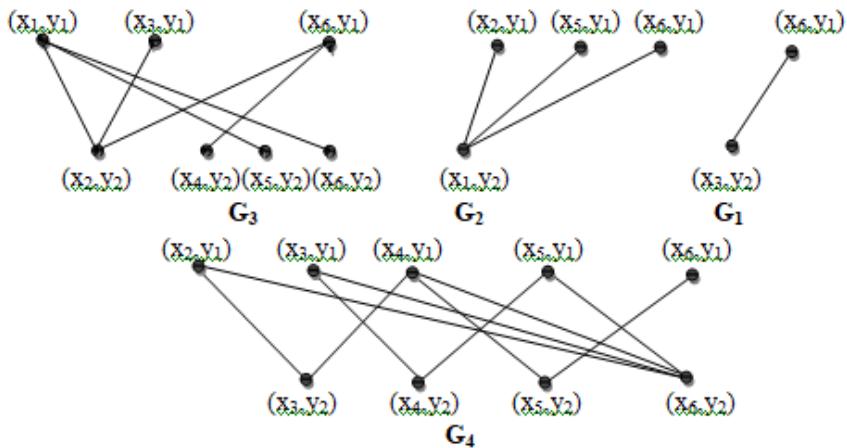


Figure 4: TD $\{G_1, G_2, G_3, G_4\}$ of $W_6 \wedge K_2$.

Theorem 2.11. For any integer m , the cycle graph $C_m \wedge K_2$ has a Triangular Decomposition $\{G_1, G_2, G_3, \dots, G_k\}$ iff there exists an integer k satisfying the following properties.

- i) $k = 4r$ or $k = 4r-1$ or $k = 4r-2$, $r \in \mathbb{N}$.
- ii) $\frac{k(k+1)(k+2)}{6} = 2m$

Proof. Let $G = C_m \wedge K_2$. Then $q(G) = 2m$. Assume $C_m \wedge K_2$ has a Triangular Decomposition $\{G_1, G_2, G_3, \dots, G_k\}$. By the definition of Triangular Decomposition, $q(G) = \binom{k+2}{3}$.

$$\text{Hence } 2m = \binom{k+2}{3} = \frac{k(k+1)(k+2)}{6}.$$

$$\Rightarrow m = \frac{k(k+1)(k+2)}{12}.$$

Since m is an integer, $k = 4r$ or $k = 4r-1$ or $k = 4r-2$, $r \in \mathbb{N}$.

Conversely assume (i) $k = 4r$ or $k = 4r-1$ or $k = 4r-2$, $r \in \mathbb{N}$ (ii) $\frac{k(k+1)(k+2)}{6} = 2m$.

Let $G = C_m \wedge K_2$. Then $q(G) = 2m$. By lemma 2.1, 2.2 and 2.3, $\frac{k(k+1)(k+2)}{6}$ can be decomposed into $\{G_1, G_2, G_3, \dots, G_k\}$. Thus G admits Triangular Decomposition.

Table 2.12: List of first 10 TD of $C_m \wedge K_2$

m	q(G)	Triangular Decomposition
5	10	G_1, G_2, G_3
10	20	G_1, G_2, G_3, G_4
28	56	$G_1, G_2, G_3, \dots, G_6$
42	84	$G_1, G_2, G_3, \dots, G_7$
60	120	$G_1, G_2, G_3, \dots, G_8$
110	220	$G_1, G_2, G_3, \dots, G_{10}$
143	286	$G_1, G_2, G_3, \dots, G_{11}$
182	364	$G_1, G_2, G_3, \dots, G_{12}$
280	560	$G_1, G_2, G_3, \dots, G_{14}$
340	680	$G_1, G_2, G_3, \dots, G_{15}$

Illustration 2.13

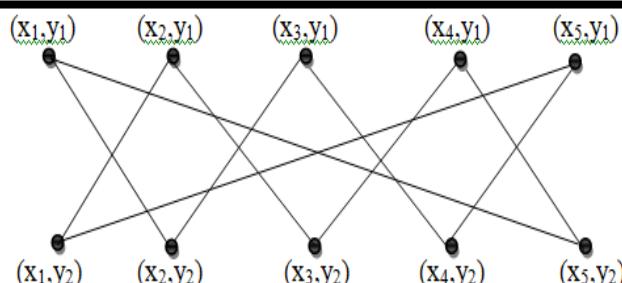


Figure 5: Tensor Product of $C_5 \wedge K_2$

Triangular Decomposition of $C_5 \wedge K_2$ as follows

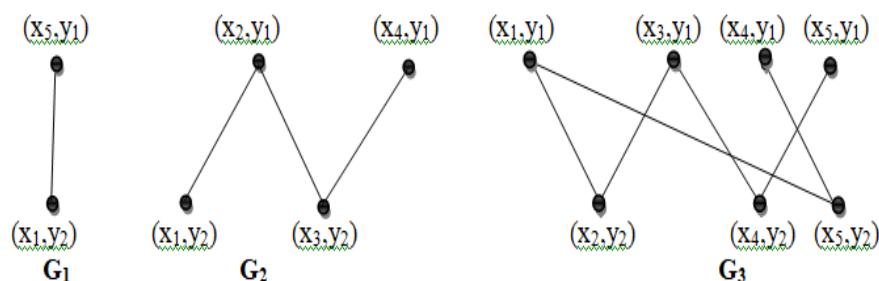


Figure 6: TD { G_1, G_2, G_3 } of $C_5 \wedge K_2$

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